

# Packet Classification Using Multidimensional Cutting



Systems & Networking

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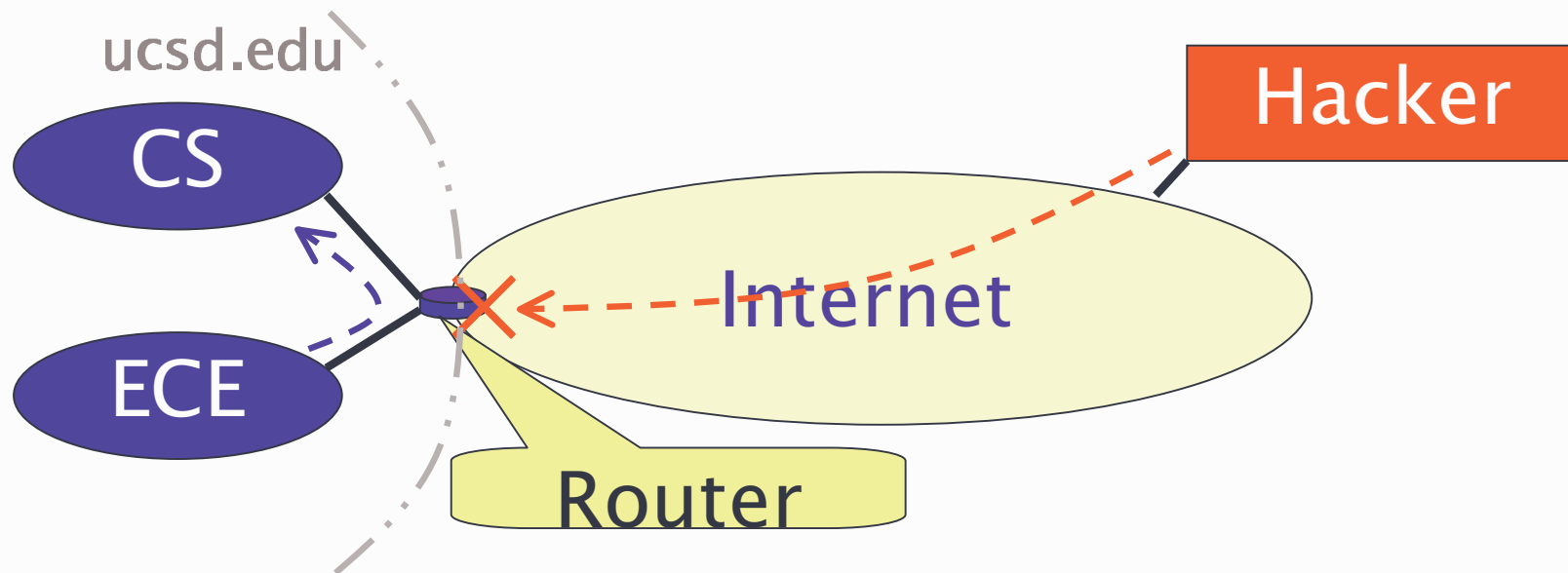
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# Packet Classification (forwarding based on multiple fields)



| Rules        | Destination | Source        | Destination Port | Action        |
|--------------|-------------|---------------|------------------|---------------|
| <i>Rule1</i> | <i>cs</i>   | <i>ece</i>    | <i>*</i>         | <i>10Gbps</i> |
| <i>Rule2</i> | <i>*</i>    | <i>hacker</i> | <i>NetBios</i>   | <i>Deny</i>   |

**Classifier** → A set of predicates (rules).

**Packet Classification** → Finding the Action associated with the highest priority rule (matching all dimensions) in the classifier.

# Rules of the Game

- Fast search **speed**  
(4–32ns /pkt throughput)
- Low **storage** requirements  
(less than several Mbits)
- **Scalability** in the **number of rules**  
(up to 100K rules)
- **Scalability** in the **number of fields**  
(five fields or more)

# Packet Classification : A Crowded Space

1998

*Bit Vector*

1999

*Grid of Tries, Crossproducing  
RFC, HiCuts*

2000

*FIS Trees*

2001

*ABV*

2003

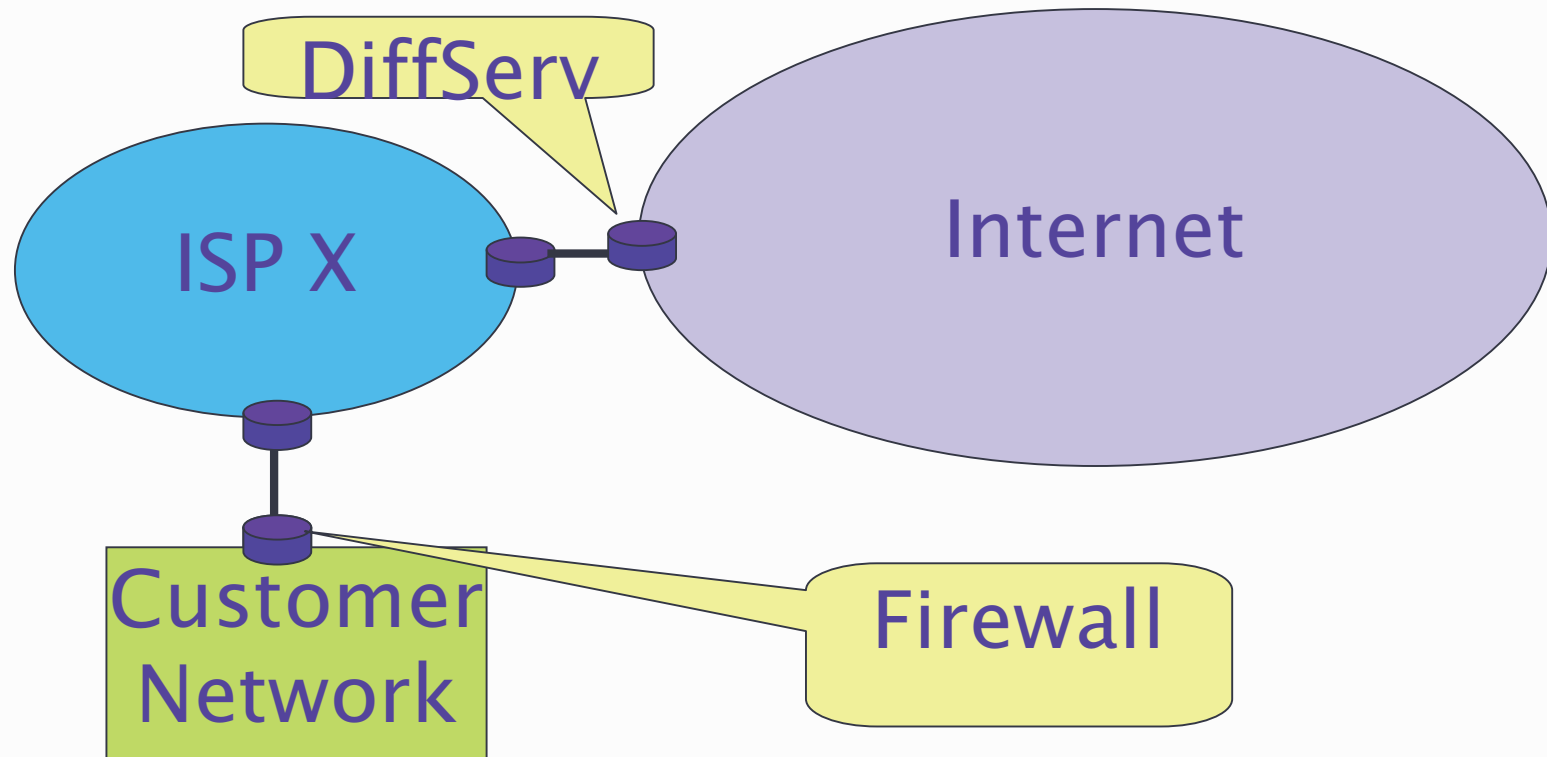
*HyperCuts (this paper)*

Why yet another paper on Packet Classification?

# Three Reasons for another solution

- A. Increasing importance of Packet classification.
- B. Inadequate performance of existing schemes:
  - CAMs
  - Algorithmic solutions
- C. Possibility of new ideas.

# A) Increasing Importance of Packet Classification



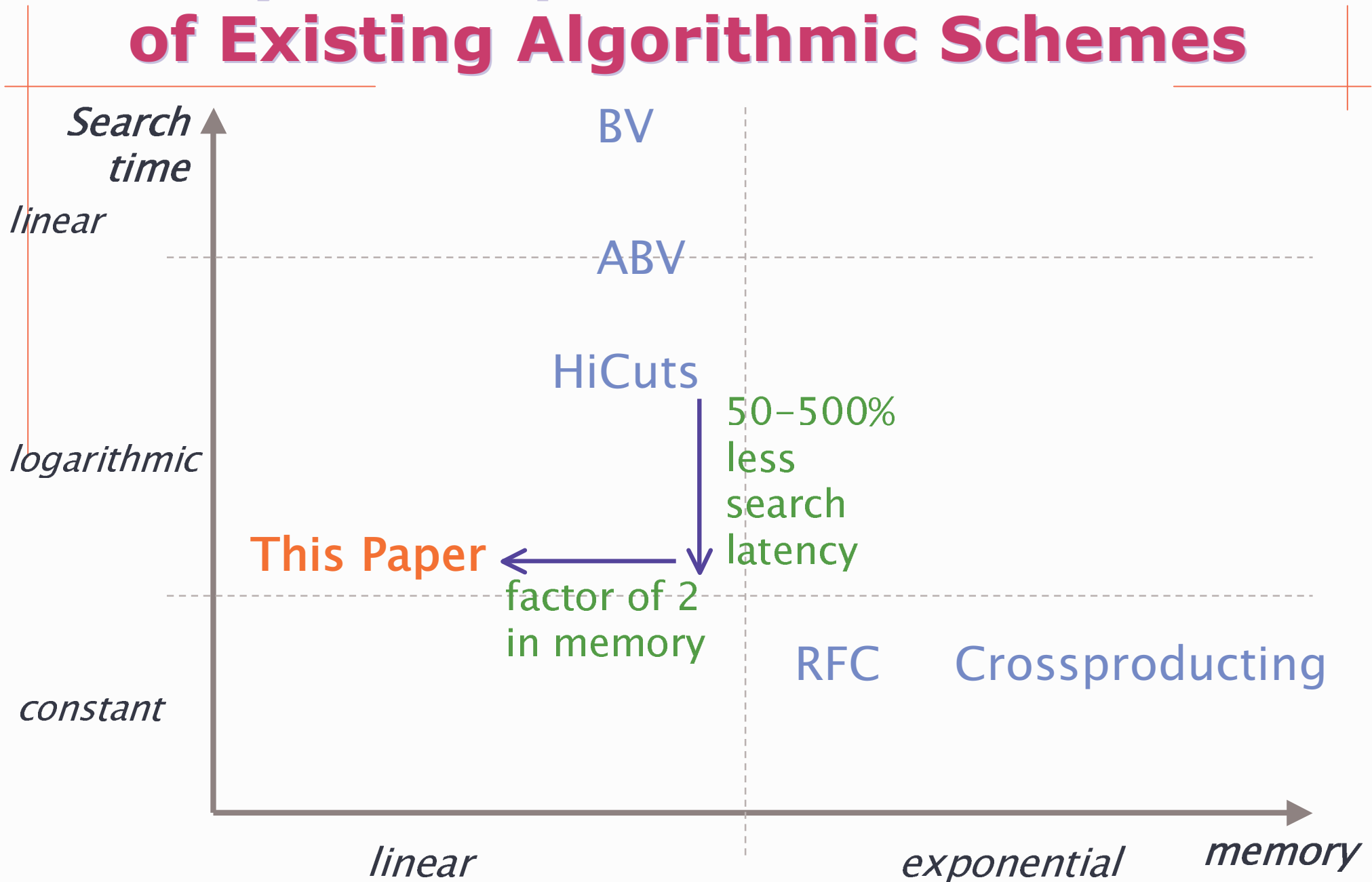
- Increased demand for new services
  - QoS
  - Security
- Increased speed
  - In 2004, 21% of edge routers will be OC-192 (10Gbps)

## B) Inadequate Performance of CAM based solutions

- Content Addressable Memory
  - Hardware Solution (using parallelism)
  - Widely used in the Industry
- Pros:
  - Low latency and high throughput
  - Simple on-chip management scheme
- Cons:
  - High power (heat!)
  - Large die size (more board space)
  - High cost (compared to SRAM based solutions)
  - All fields must be expressed into a prefix format

**An algorithmic solution may be a contender!**

# B) Inadequate Performance of Existing Algorithmic Schemes





## C) Possibility of New Ideas

- Main Idea:
  - Increasing degrees of freedom involved in decision tree approaches to classification, by using **hypercubes** to partition the search space instead of **hyperplanes**.

# Outline

1. Introduction

➔ 2. Geometric View of Packet Classification

3. Basic Decision Tree Approaches

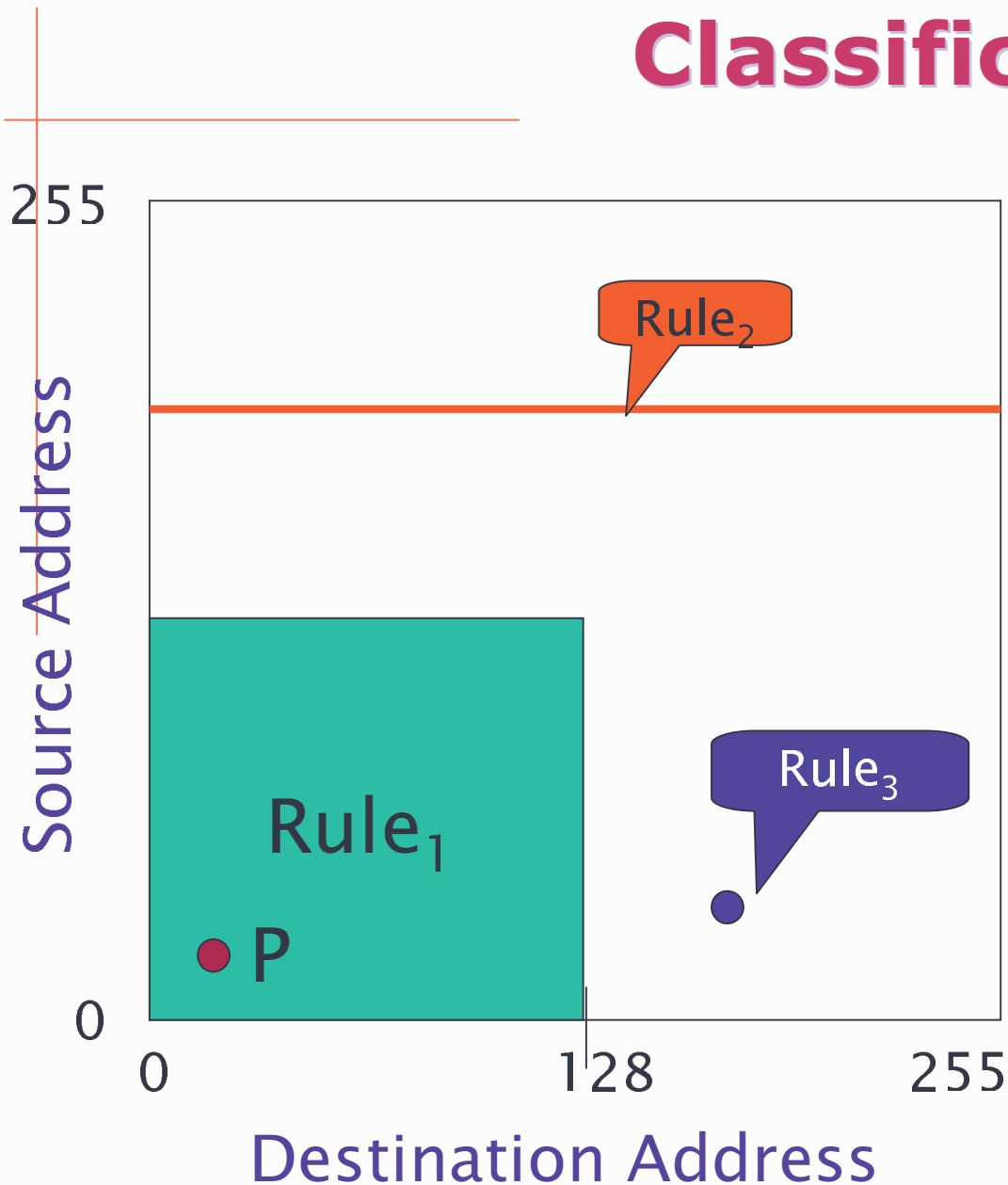
4. Basic HyperCuts

5. HyperCuts Optimizations

6. Experimental Results

7. Conclusion

# Geometric View of Packet Classification



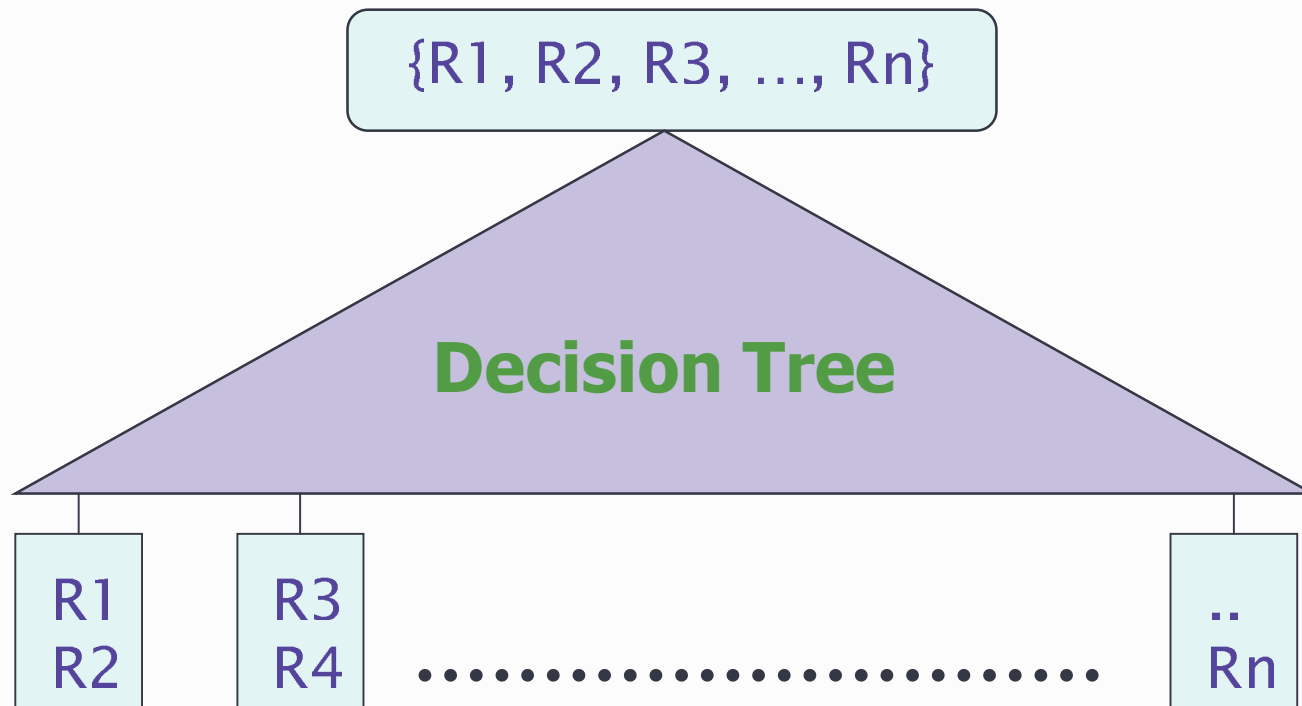
| Rules             | Source  | Destination |
|-------------------|---------|-------------|
| Rule <sub>1</sub> | 0 - 127 | 0 - 127     |
| Rule <sub>2</sub> | 192     | 0 - 255     |
| Rule <sub>3</sub> | 32      | 160         |

Prefixes represented as ranges

# Outline

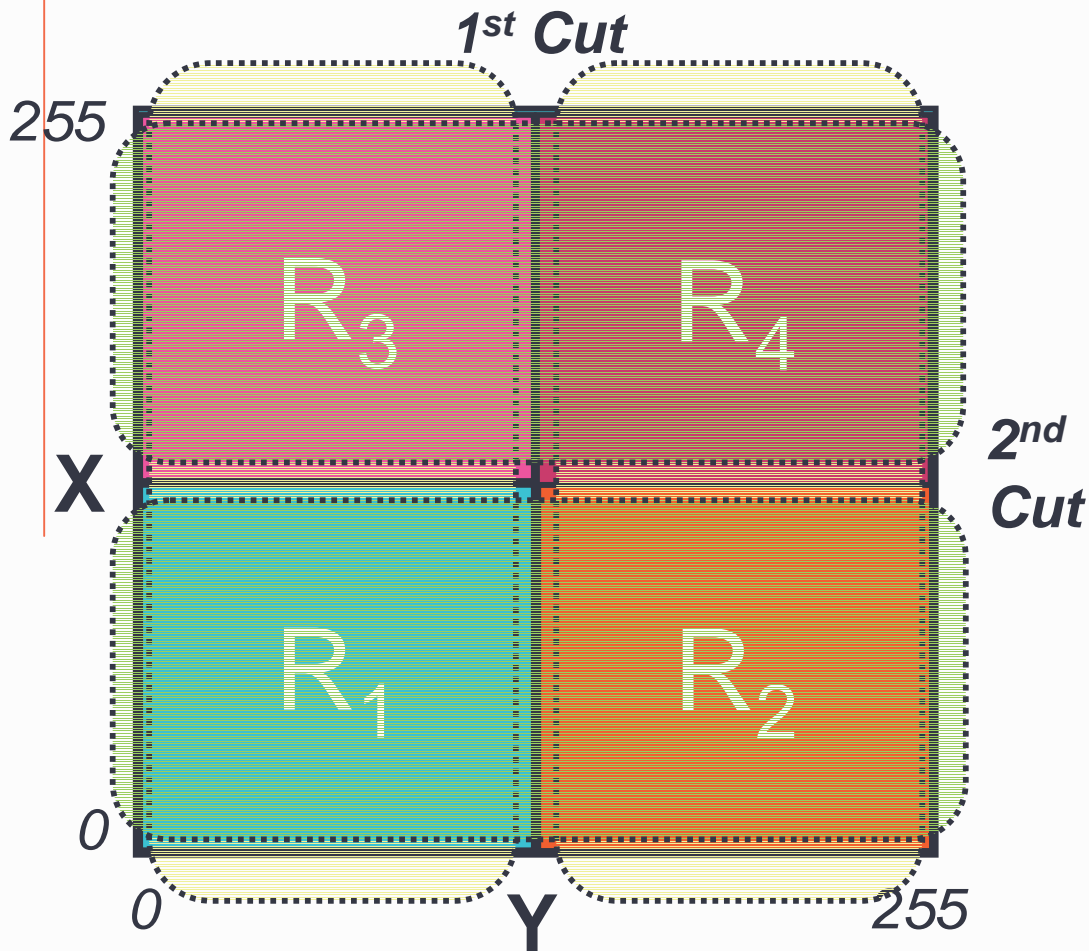
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# Decision Tree Based Classification

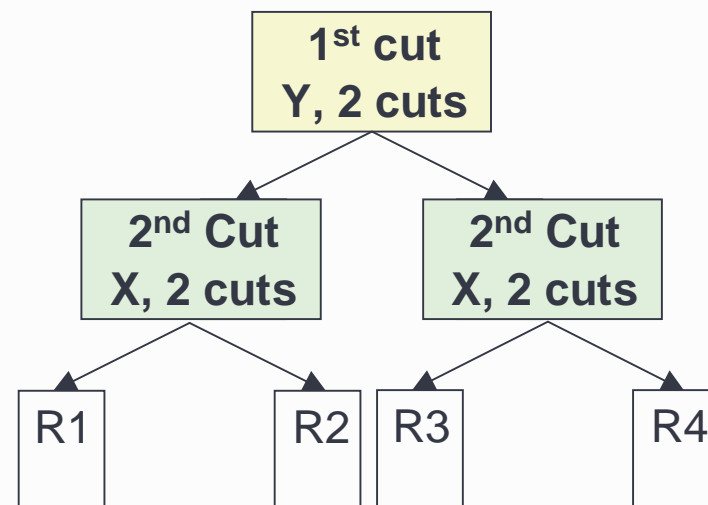


Pioneered by Woo and Gupta-McKeown

# HiCuts: Using single-dimension cutting




|                |    |    |
|----------------|----|----|
| R <sub>1</sub> | 0* | 0* |
| R <sub>2</sub> | 0* | 1* |
| R <sub>3</sub> | 1* | 0* |
| R <sub>4</sub> | 1* | 1* |

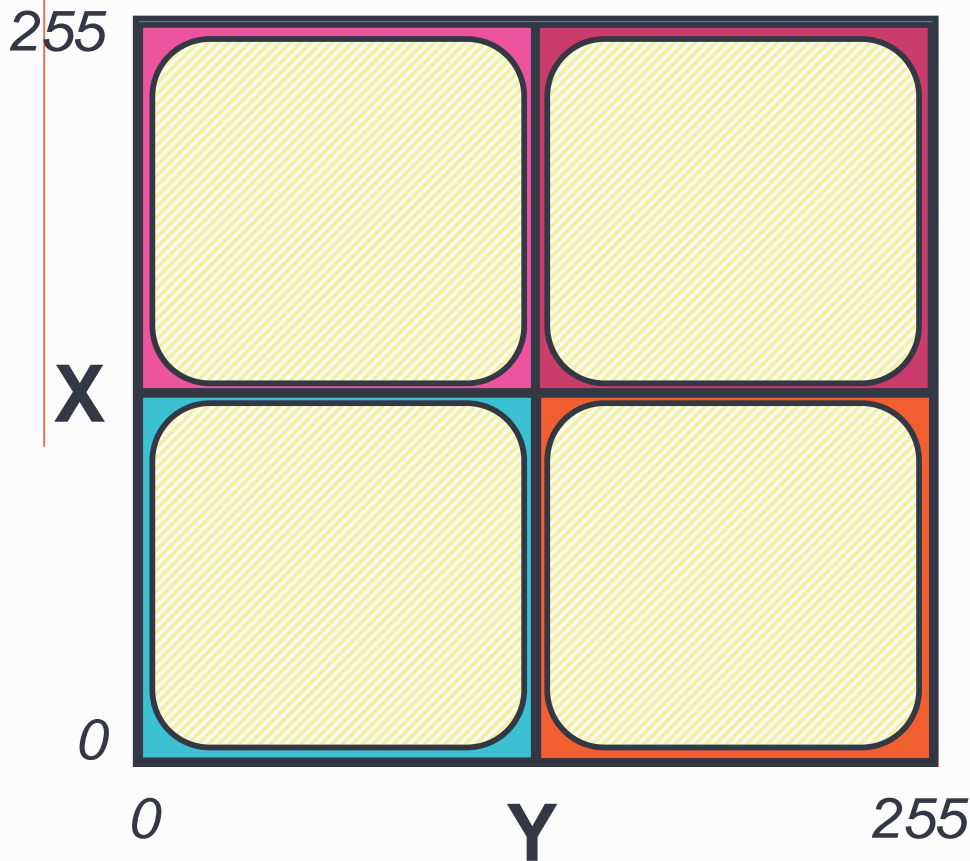


Gupta-McKeown 99

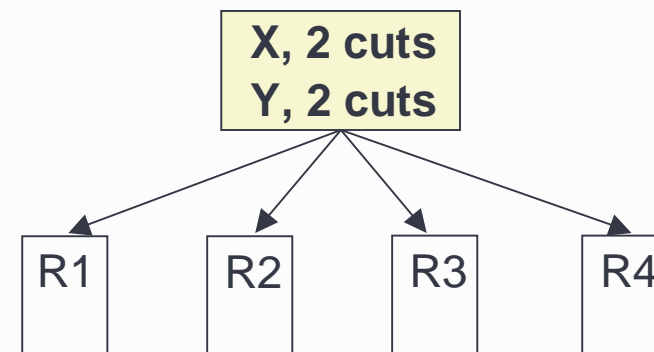
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# Using multidimensional cutting



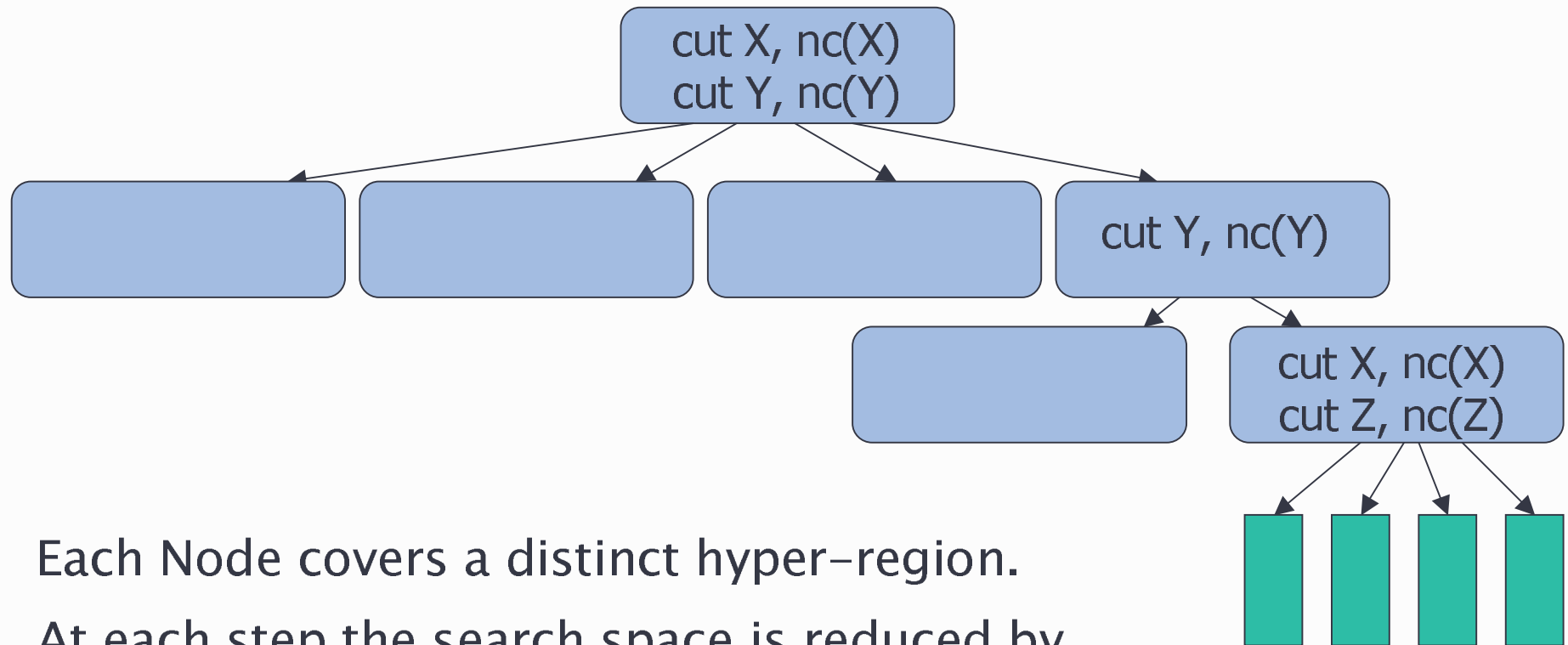
|                |    |    |
|----------------|----|----|
| R <sub>1</sub> | 0* | 0* |
| R <sub>2</sub> | 0* | 1* |
| R <sub>3</sub> | 1* | 0* |
| R <sub>4</sub> | 1* | 1* |



Cuts are **equal size ranges** on each dimension, for easy array indexing. The number of cuts in each dimension may be **different**.



# A HyperCuts Decision Tree

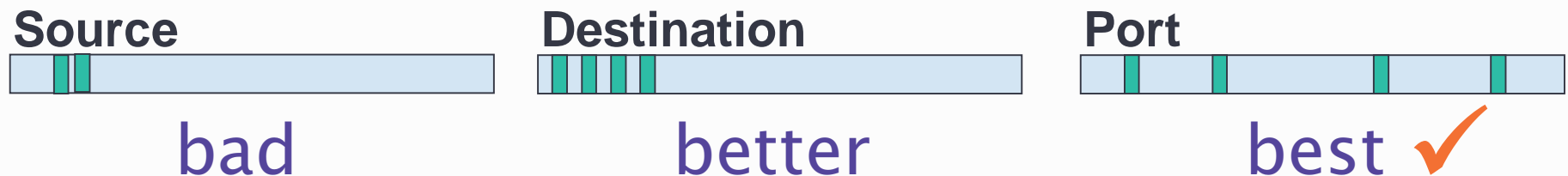


- Each Node covers a distinct hyper-region.
- At each step the search space is reduced by cutting a node (across  $k$ -dimensions).
- All child-nodes of the same parent cover non-overlapping hyper-regions of same size.
- Leaf-Nodes have a small number of rules represented in a list.

# Building the HyperCuts decision tree

## Step 1: Selecting the Dimensions

- Challenge:
  - To pick the dimensions which will lead to the most uniform distribution of the rules when the node is cut into sub-nodes.
- Idea:
  - Pick dimensions with highest entropy.



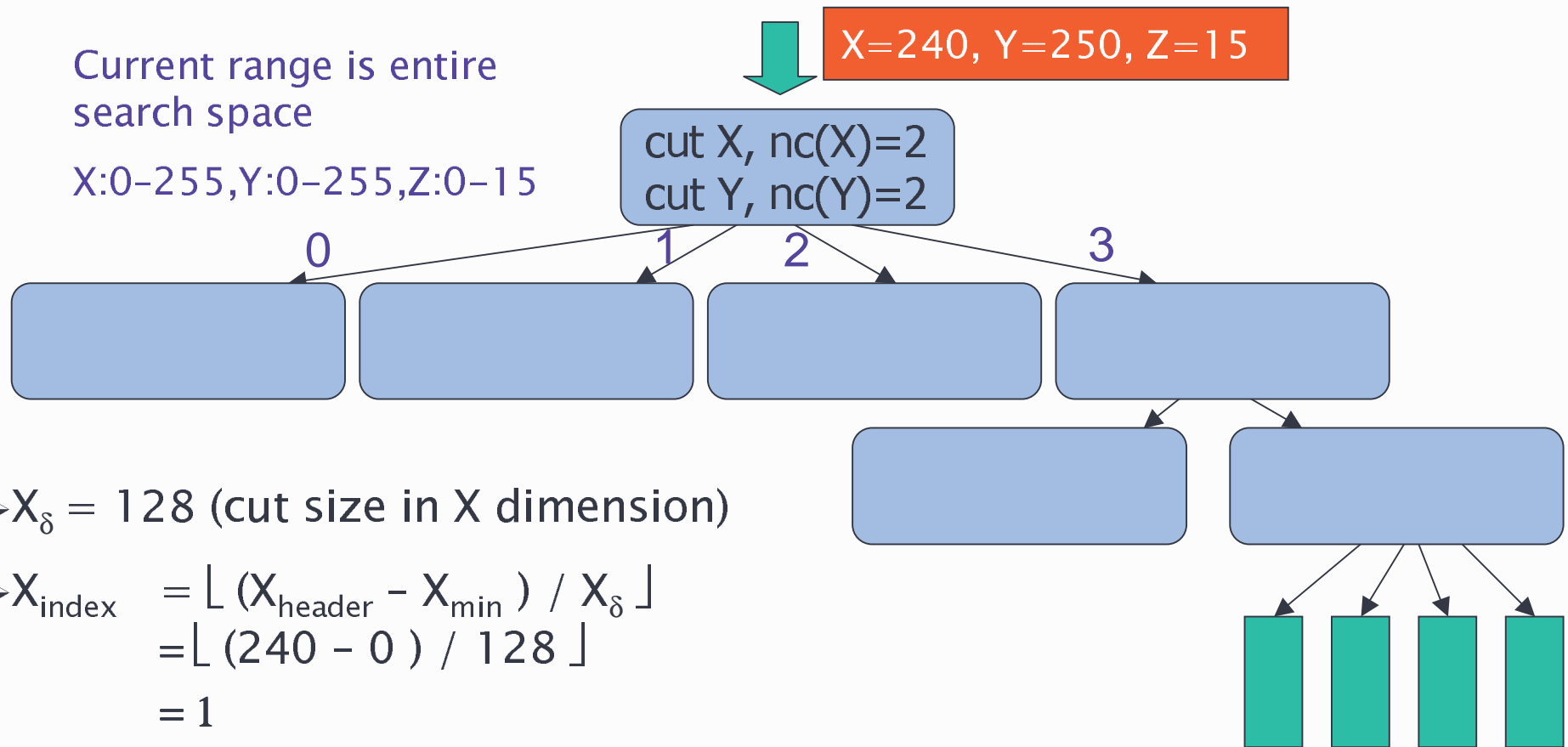
Recall: cuts are equal size ranges for easy array indexing!

# Building the HyperCuts decision tree

## Step 2: Selecting the # of cuts

- Goal 1: Minimize search time while keeping space roughly linear
- Strategy 1: Look for multi-dimensional cut that:
  - Minimizes number of rules allocated to any child node
  - Maximum number of Children (cuts) allocated to a node are limited by (space factor \*  $\sqrt{\text{\#rules in node}}$ ).
- Goal 2: Avoid exponential time to create a good decision tree
- Strategy 2: Use a greedy strategy which:
  - Determines the optimal cut in each dimension
  - Considers only combinations of these locally optimal cuts

# Search algorithm for a HyperCuts decision tree



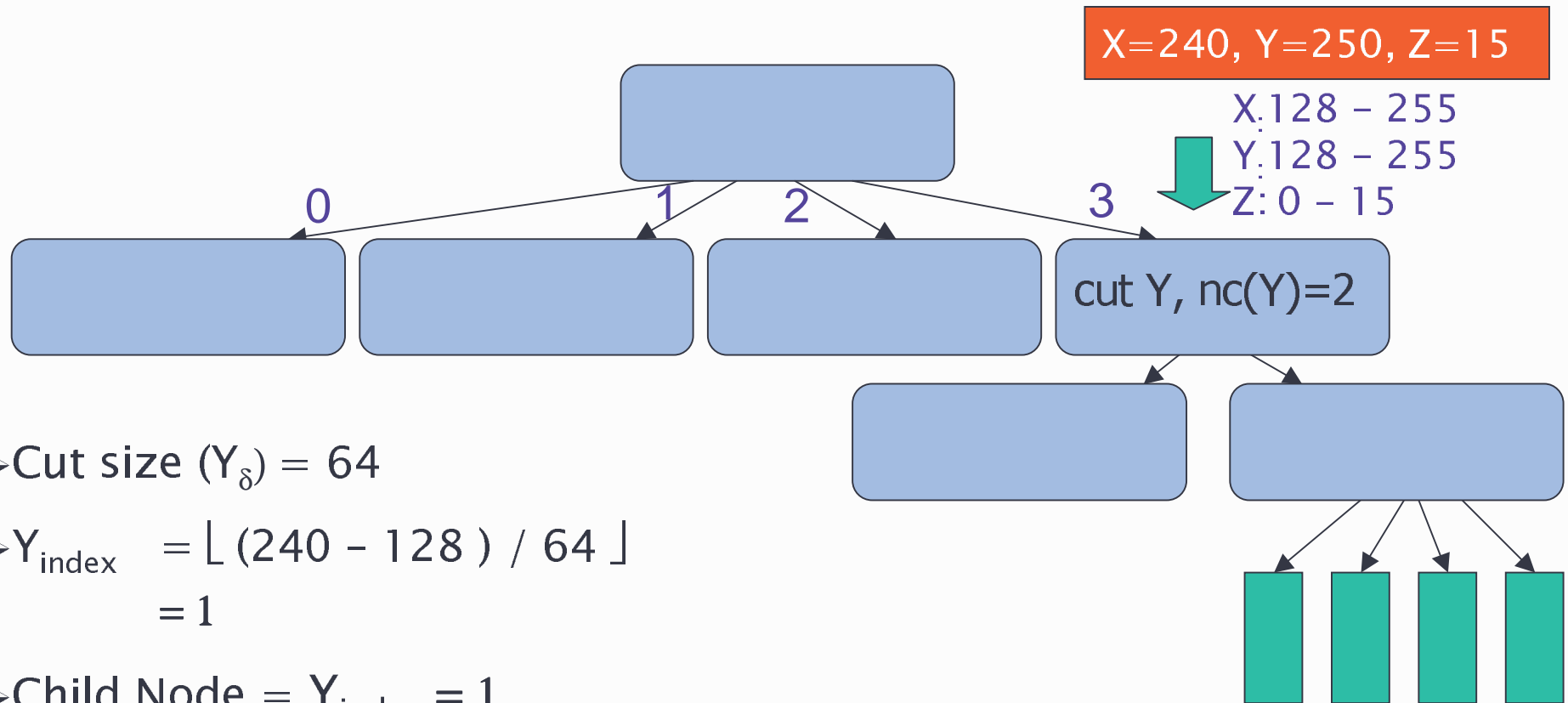
➤  $X_{\delta} = 128$  (cut size in X dimension)

$$\begin{aligned} \text{➤ } X_{\text{index}} &= \lfloor (X_{\text{header}} - X_{\text{min}}) / X_{\delta} \rfloor \\ &= \lfloor (240 - 0) / 128 \rfloor \\ &= 1 \end{aligned}$$

$$\text{➤ } Y_{\text{index}} = \lfloor (250 - 0) / 128 \rfloor = 1$$

$$\begin{aligned} \text{➤ Child Node} &= Y_{\text{index}} * nc(Y) + X_{\text{index}} \\ &= (1 * 2) + 1 = 3 \end{aligned}$$

# Search algorithm for a HyperCuts decision tree

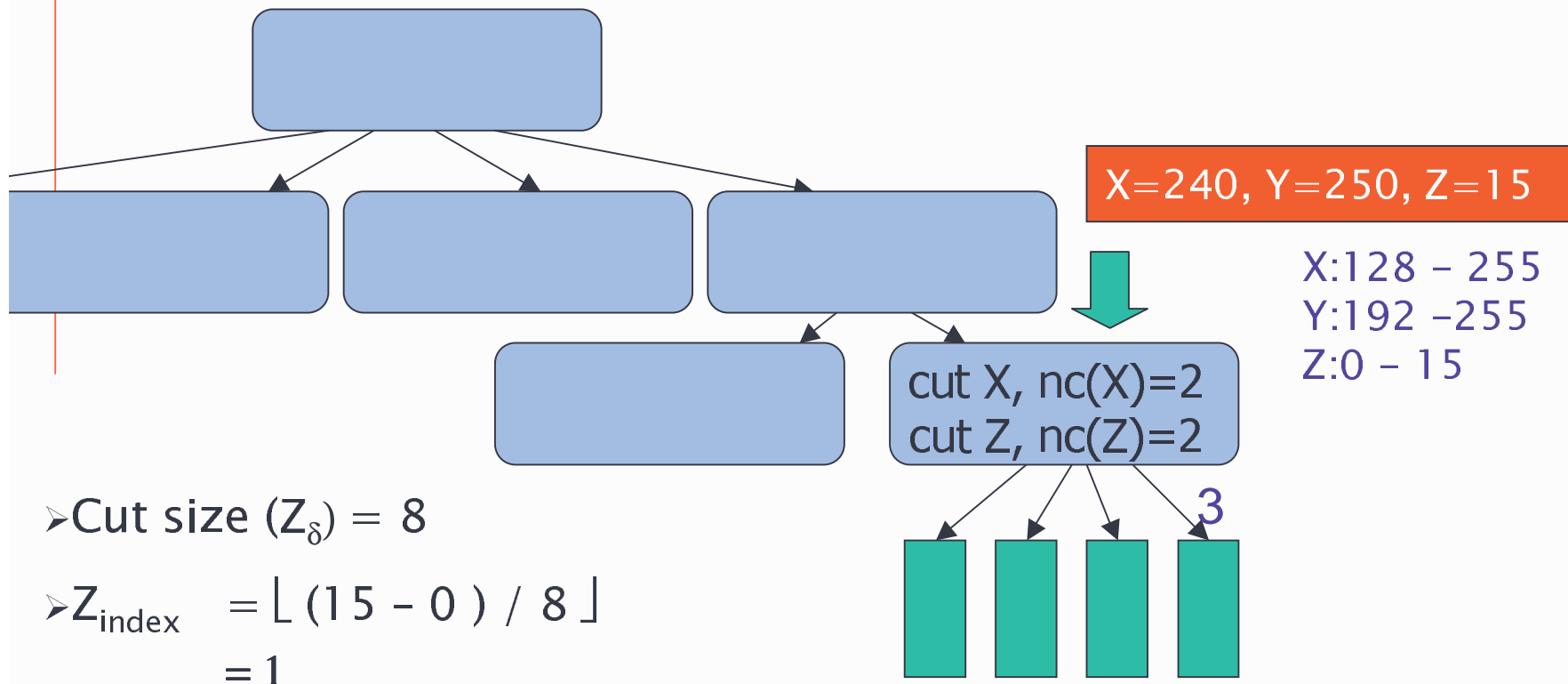


➤ Cut size ( $Y_\delta$ ) = 64

➤  $Y_{\text{index}} = \lfloor (240 - 128) / 64 \rfloor$   
= 1

➤ Child Node =  $Y_{\text{index}} = 1$

# Search algorithm for a HyperCuts decision tree



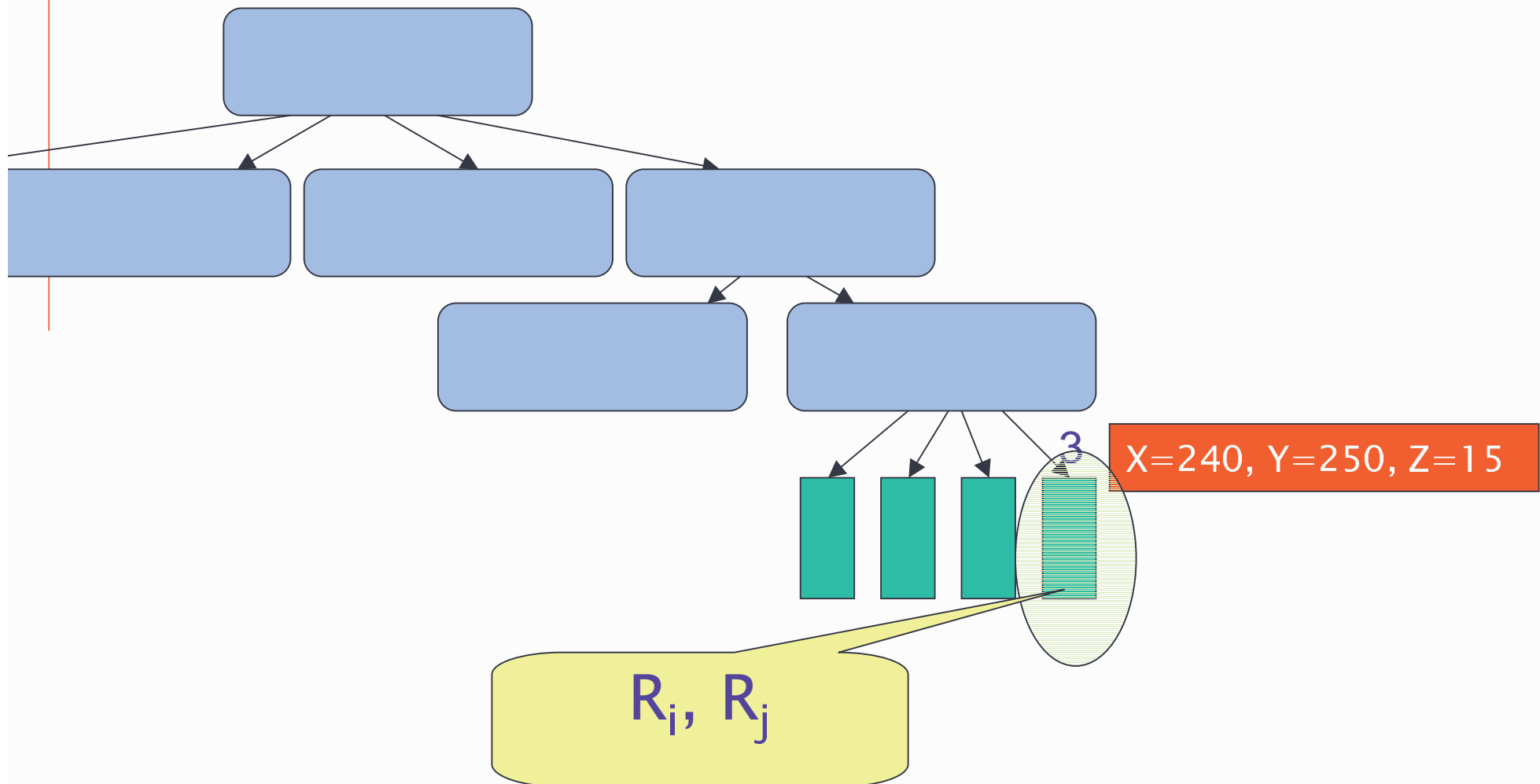
➤ Cut size ( $Z_\delta$ ) = 8

➤  $Z_{\text{index}} = \lfloor (15 - 0) / 8 \rfloor = 1$


➤  $X_{\text{index}} = \lfloor (240 - 128) / 64 \rfloor = 1$

➤ Child Node =  $Z_{\text{index}} * nc(Z) + X_{\text{index}}$   
 $= (1 * 2) + 1 = 3$

# Search algorithm for a HyperCuts decision tree



# Outline

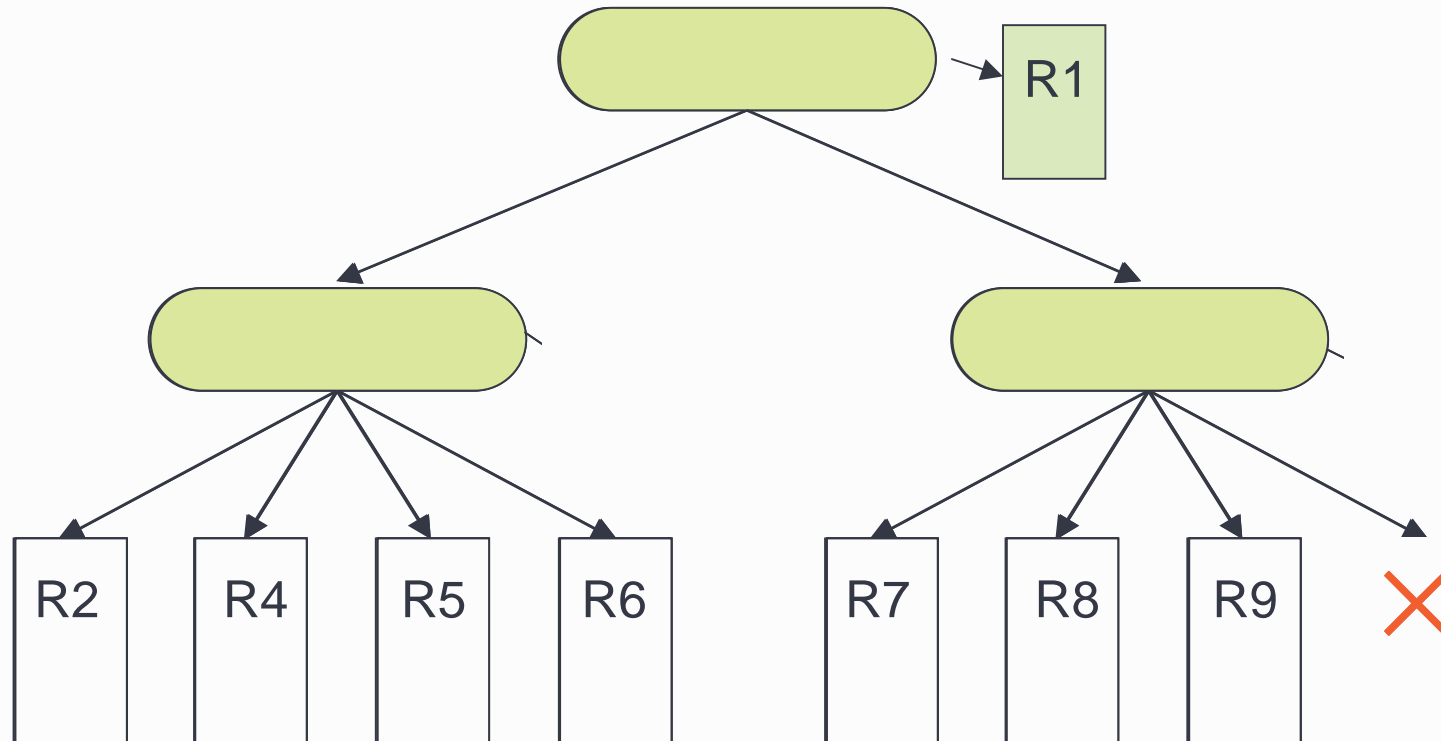
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# Optimizations for Space Reduction


- Two sources of memory wastage in basic HyperCuts
  - Space consumed by multidimensional arrays.  
Solutions: Node merging, Region compaction
  - Space consumed by *replicated rules*.  
Solutions: Eliminate Rule overlap, **Rule Pushing**

# Rule Pushing



- Rule R1 exists in all child-nodes
- Push-up rule R1 to parent node
- Wild carded rules often get replicated like this.

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# Evaluation Methodology

## ➤ Metrics:

- Worst case search time in number of memory accesses
- Memory size

## ➤ Real & Synthetic Classifiers:

- Core routers (real from multiple Tier-1 ISPs)
- Edge routers
- Firewalls

## Notes:

Each rule in the classifiers is a 5 Tuple:

Source Prefix, Destination Prefix, Source Port, Destination Port, Protocol

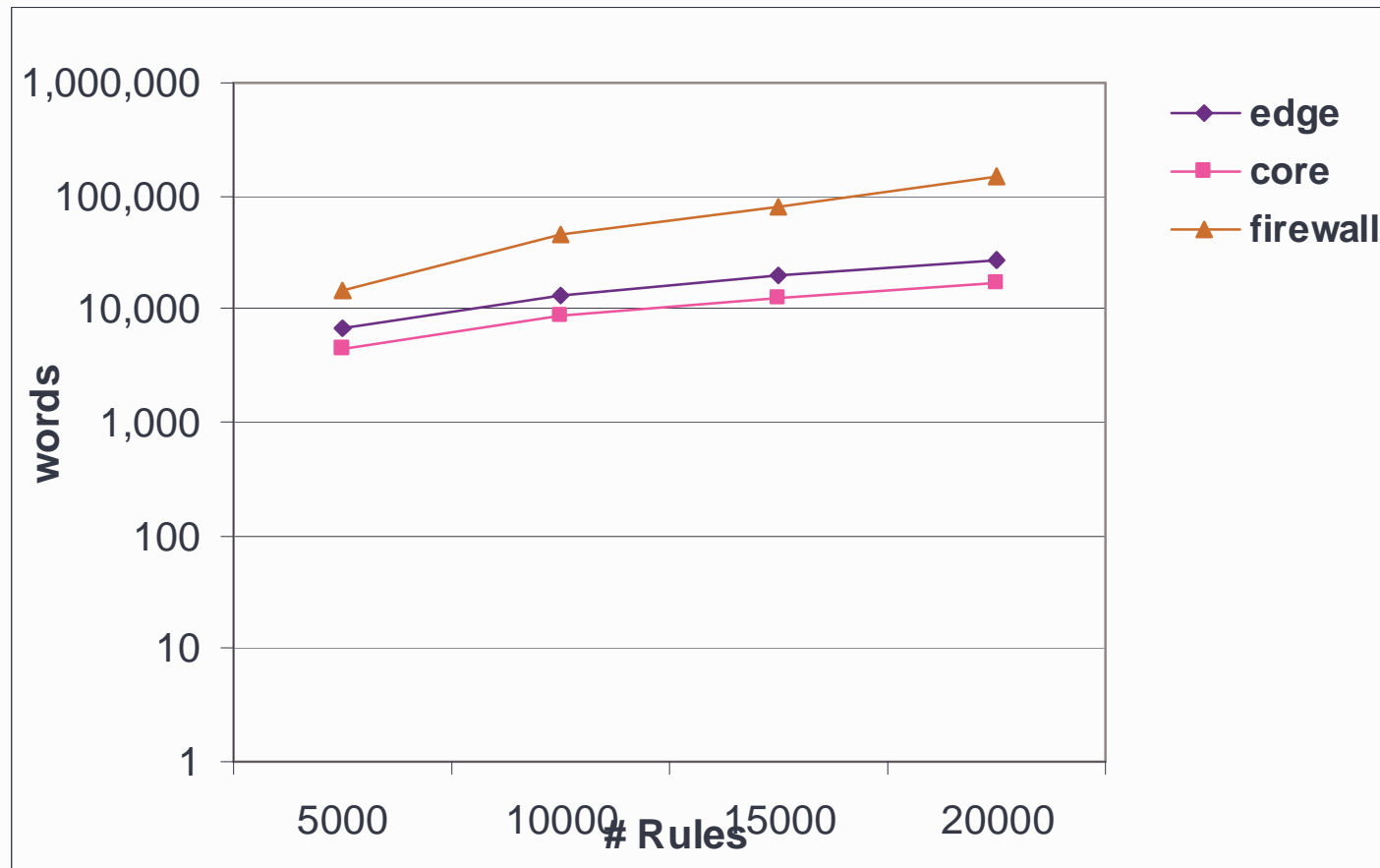
# Evaluation

## Real Classifiers

- HyperCuts optimized for memory has **50–500% better search time** than HiCuts optimized for speed.
- HyperCuts optimized for speed uses **2 to 10 times less memory** than HiCuts optimized for memory.
- Compared with other algorithms (e.g. RFC) for a database of 2800 rules HyperCuts uses **30 times less memory space**, while the search speed decreases only by a factor of 50%.

# Evaluation

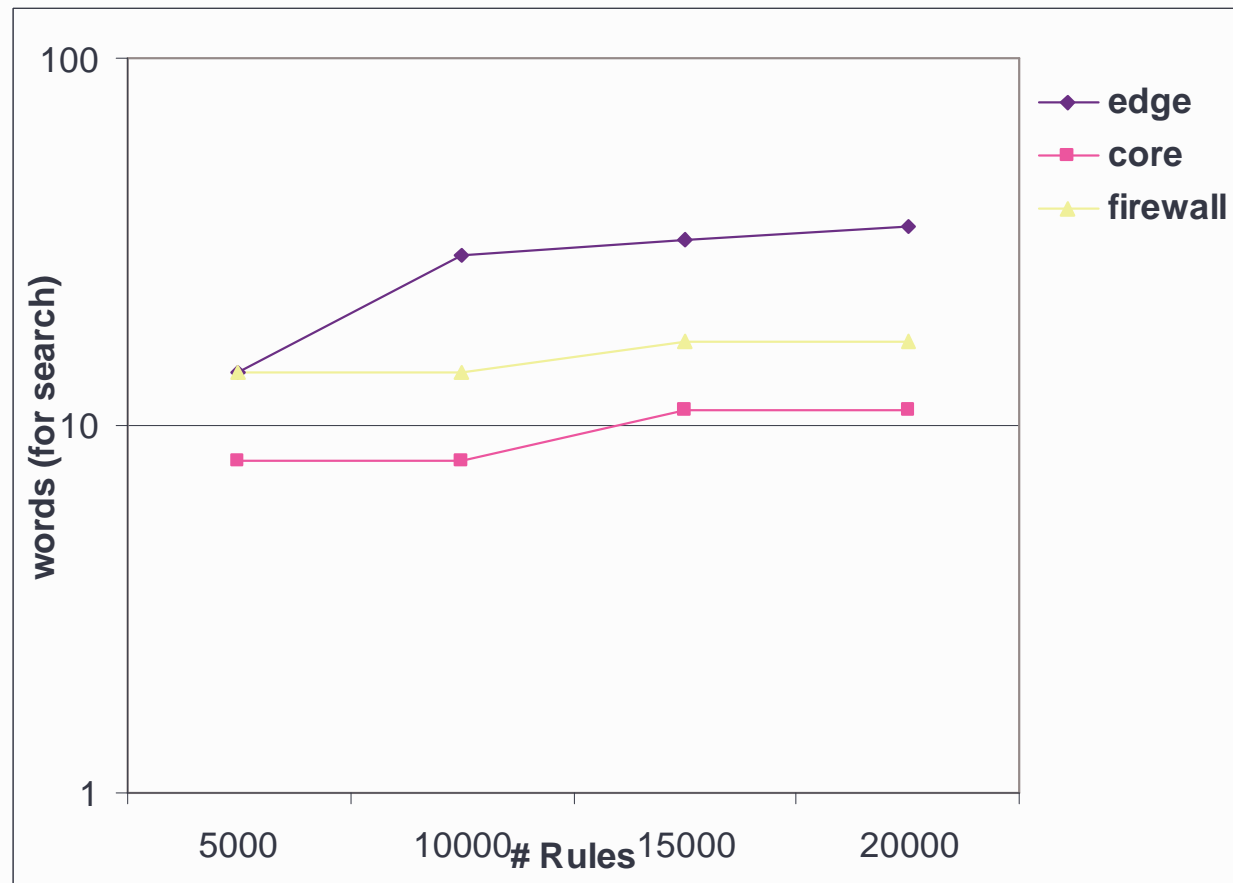
## Synthetic classifiers (memory)



- Memory utilization grows linearly with increase in number of rules

# Evaluation

## Synthetic Classifiers (search)



- Search time does not grow worse than logarithmically

# A word of caution

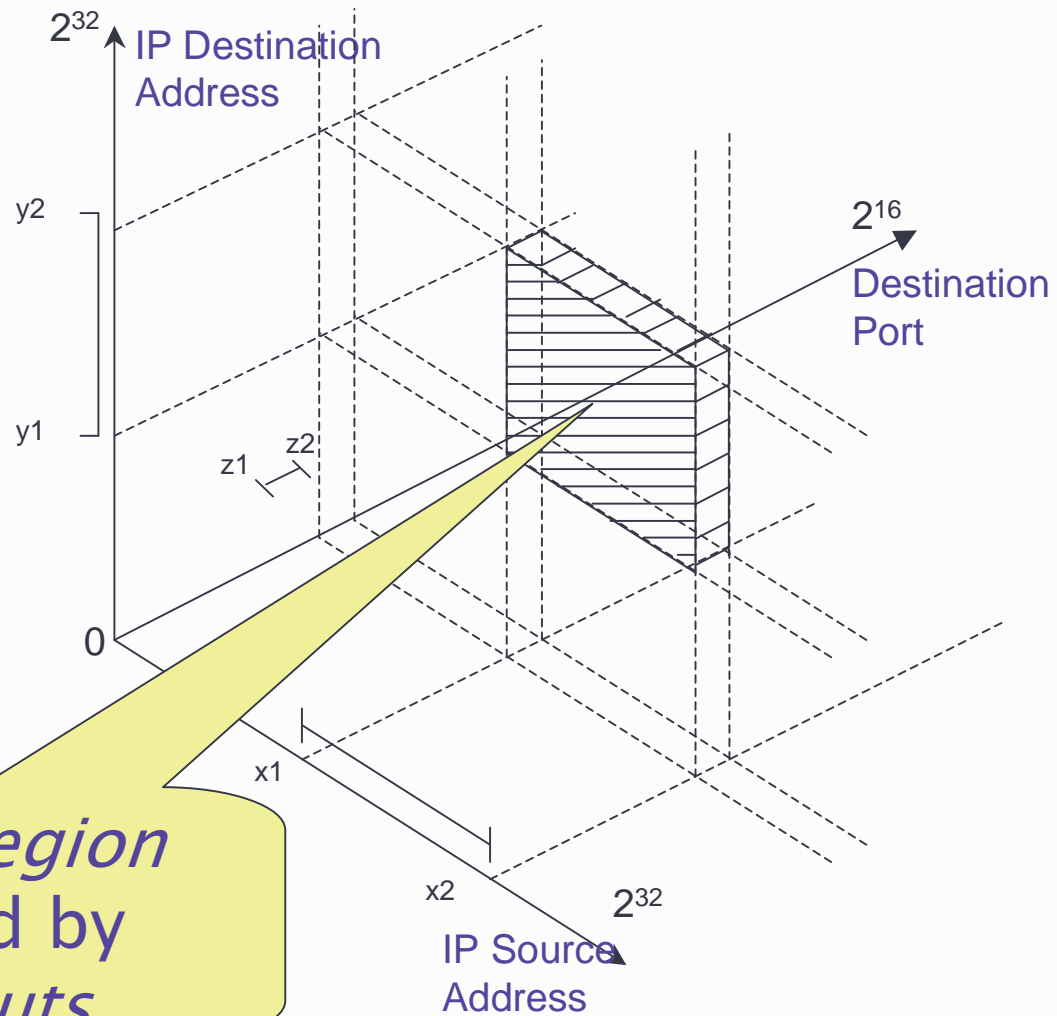
- Classifier characteristics differ between locations and between ISPs (Firewall, Edge, Core Router)
- Cutting across multiple dimensions in each step may not be a good idea:
  - Lose flexibility of adaptive decisions
- For 2-d classifiers HyperCuts degenerates to HiCuts for best performance (i.e. select at most 1 dimension at every step)



# Conclusion

- HyperCuts has linear space complexity and provides a latency that is at most logarithmic in the number of rules on real classifiers that we studied.
- The throughput of the algorithm can be improved by **pipelining** based on the depth of the tree.
- Based on initial evaluation, It seems that HyperCuts can be a practical contender compared to CAM based solutions.
- Future Direction:  
We have designed a pipeline architecture for hardware implementation of the algorithm, which we are evaluating.

# Questions ?



*A HyperRegion  
produced by  
HyperCuts*

# Decision Tree Based Algorithms

## ➤ Idea:

- build a decision tree based on local optimization decisions at each node

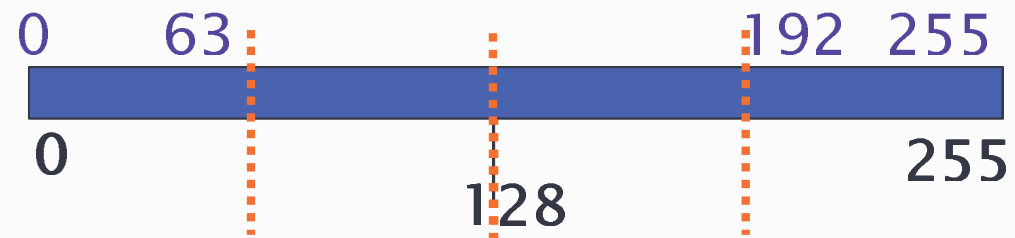
## ➤ Pros:

- Tree can be of relatively small height
- Easy to pipeline

## ➤ Cons:

- Difficult to predict the performance
- Utilizing fancy heuristics and optimizations may
  - Increase search latency
  - Increase complexity of incremental updates.

# What is a Cut?



4 cuts, cut size= 64